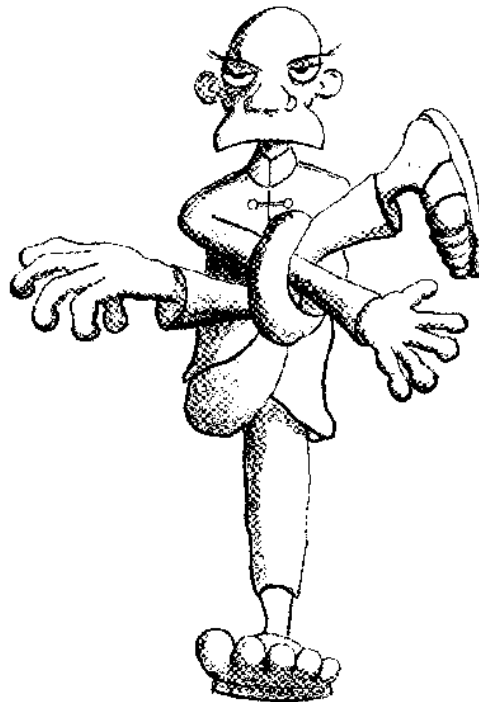


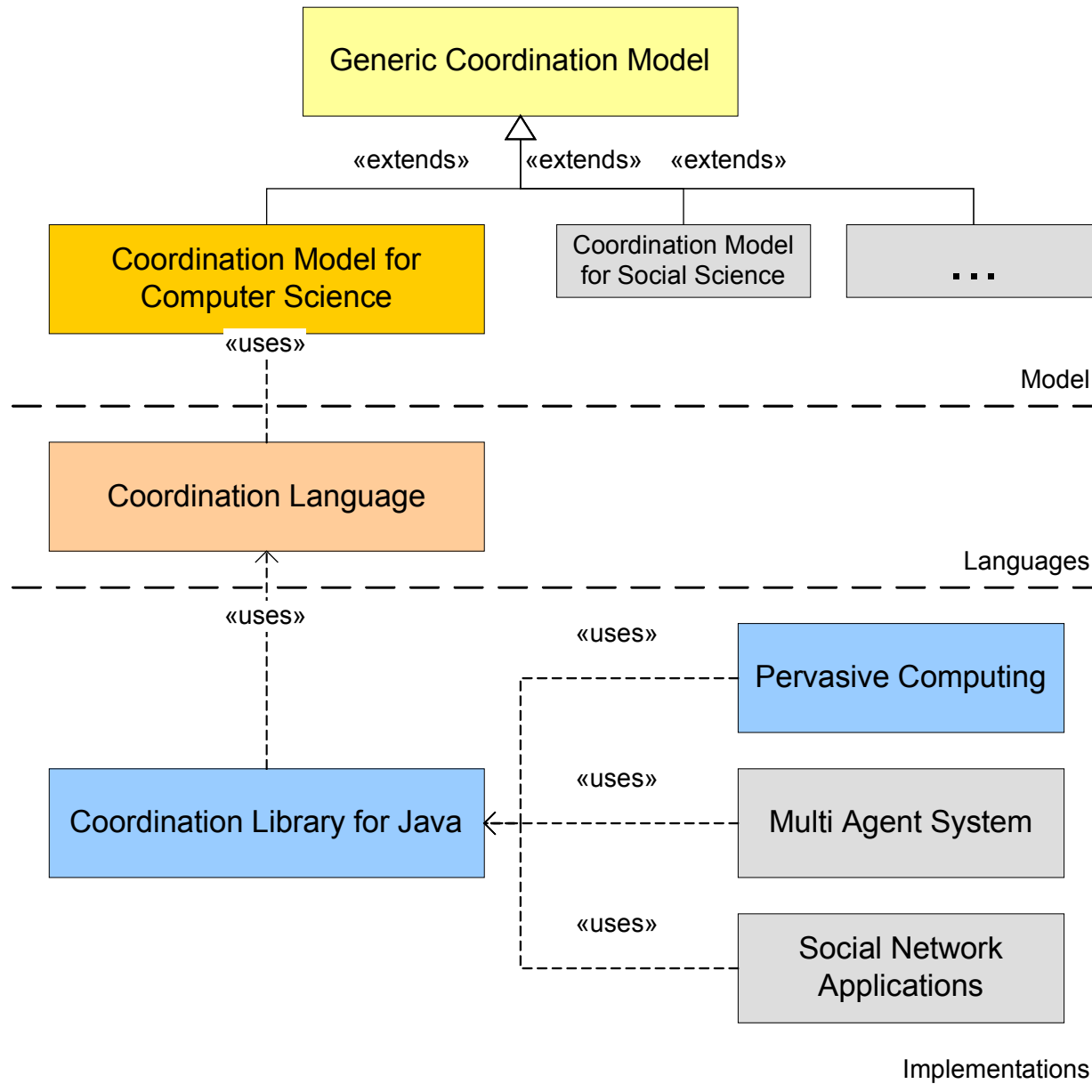
Coordination Language

(illustrated with Java)

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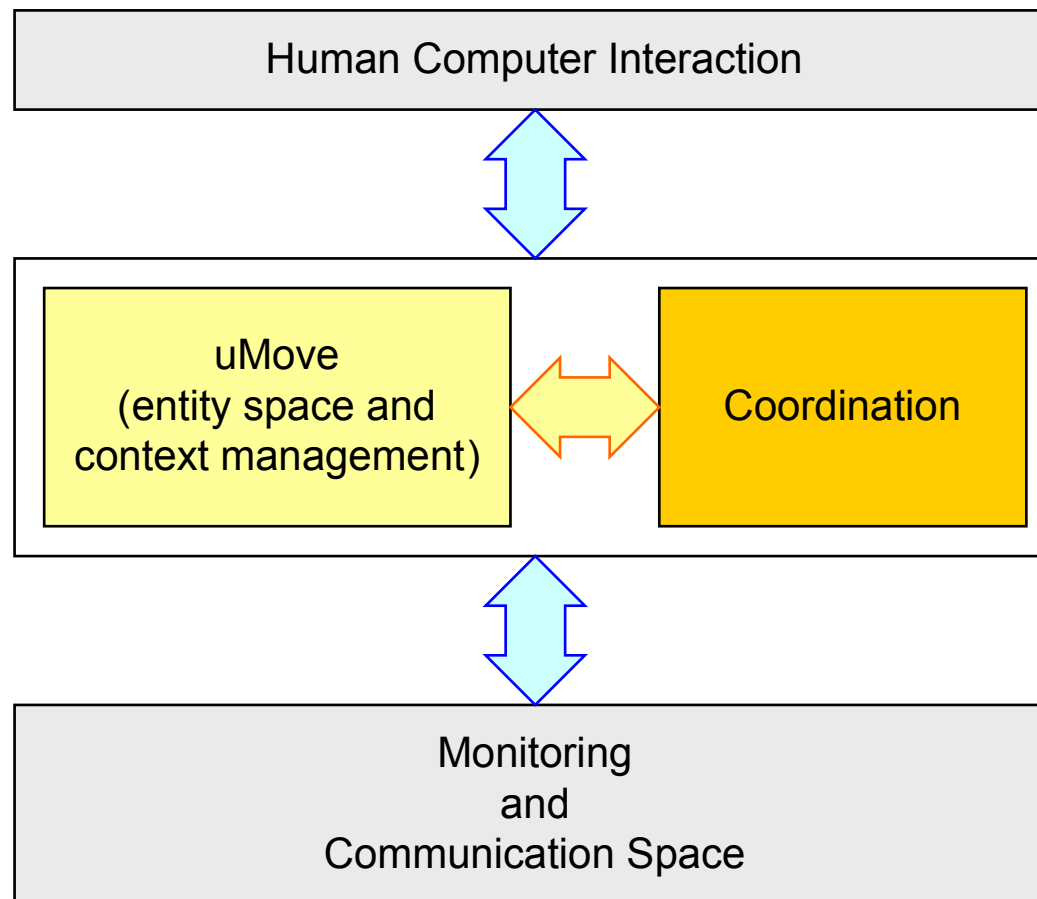


Overview

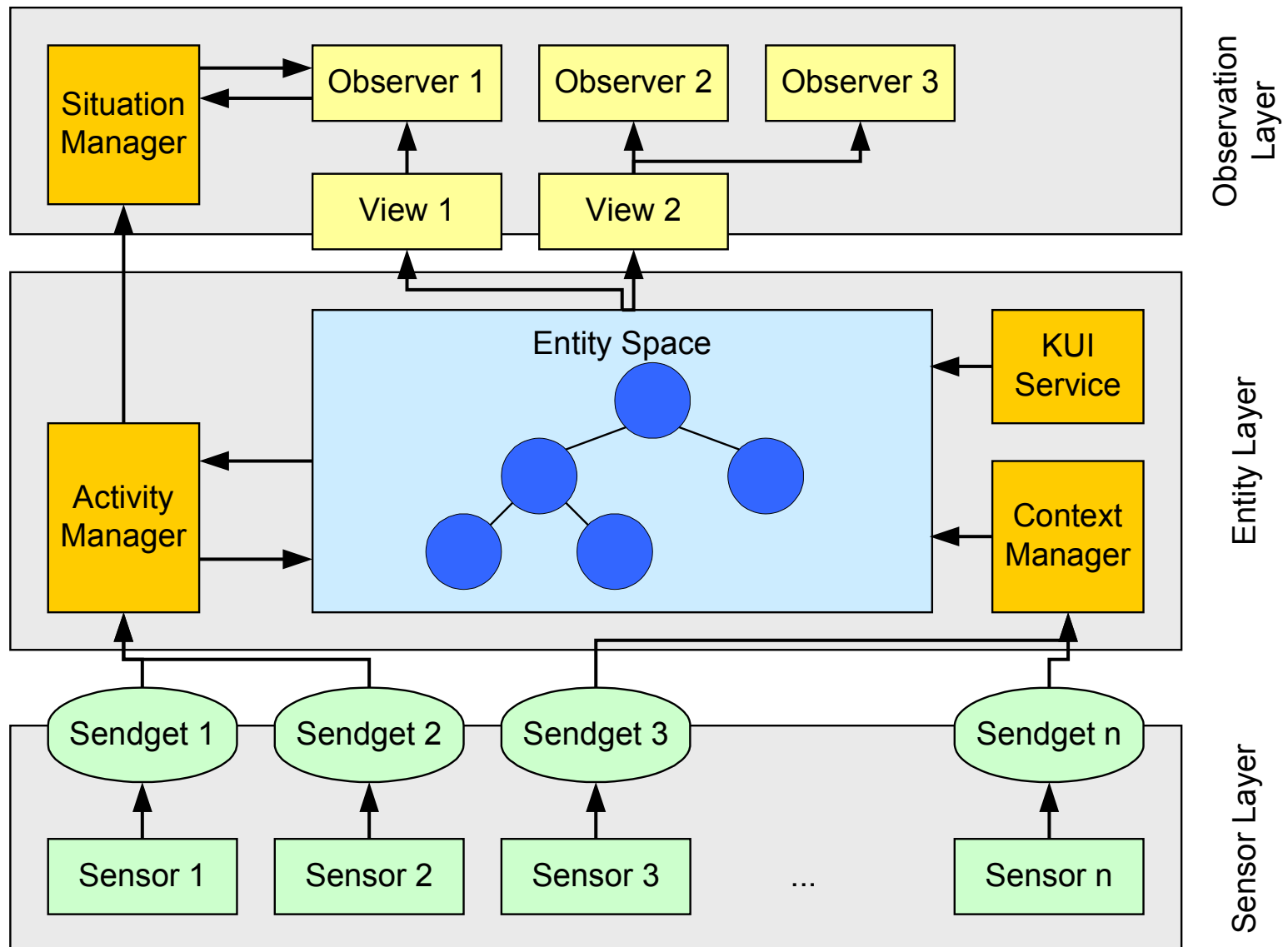


Architecture of Middleware for Pervasive Computing

- Java Package extending the computational part of the middleware by coordination



uMove: Representation of the World



Entities

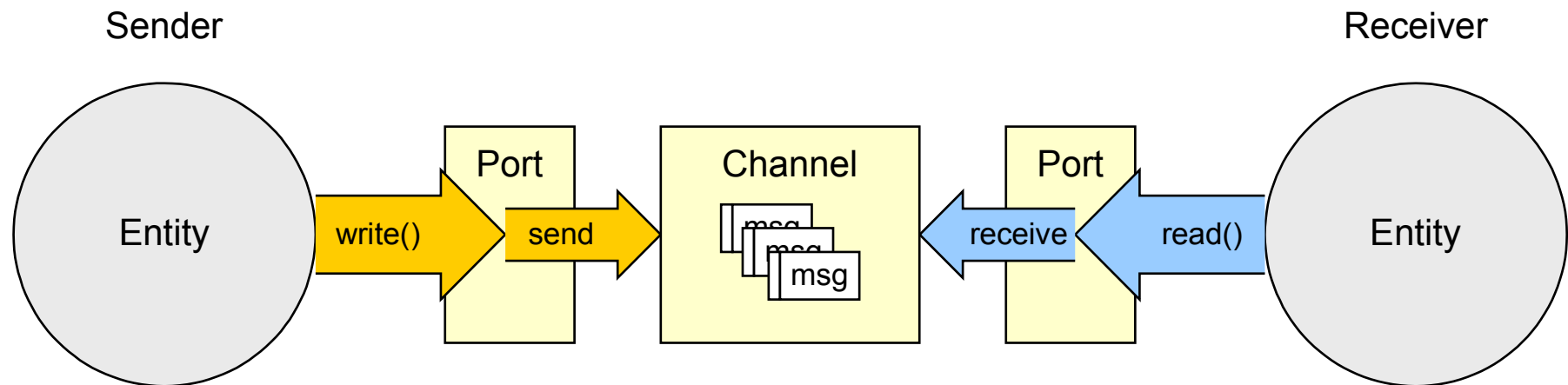
- Entities are described by
 - Identity: unique ID (UUID), Name, RFID Tag, etc.
 - Context: Sensor values, location
 - Motion and Activity
- There exist different types:
 - Actor: observed entities of the physical world
 - Observer: Software observer looking at a part of the world
 - Viewer: Software filter used
 - Sendget: Sensor entity
 - Group: virtual entity used to group physical entities

Ports (I)

- Ports are described by a port descriptor:
 - ID and Name
 - Orientation (input, output or inout)
 - Synchronisation type (asynchron or synchron)
 - Delivery Protocol (FIFO, causal order, total order, ...)
 - Public flag (visible for others)
 - Address

Ports (II)

- An entity can access the communication channel only through ports
 - **write**: The sender puts a message into the channel. The port must have output access.
 - **read**: The receiver reads a message from the channel. The port must have input access.



Ports (III)

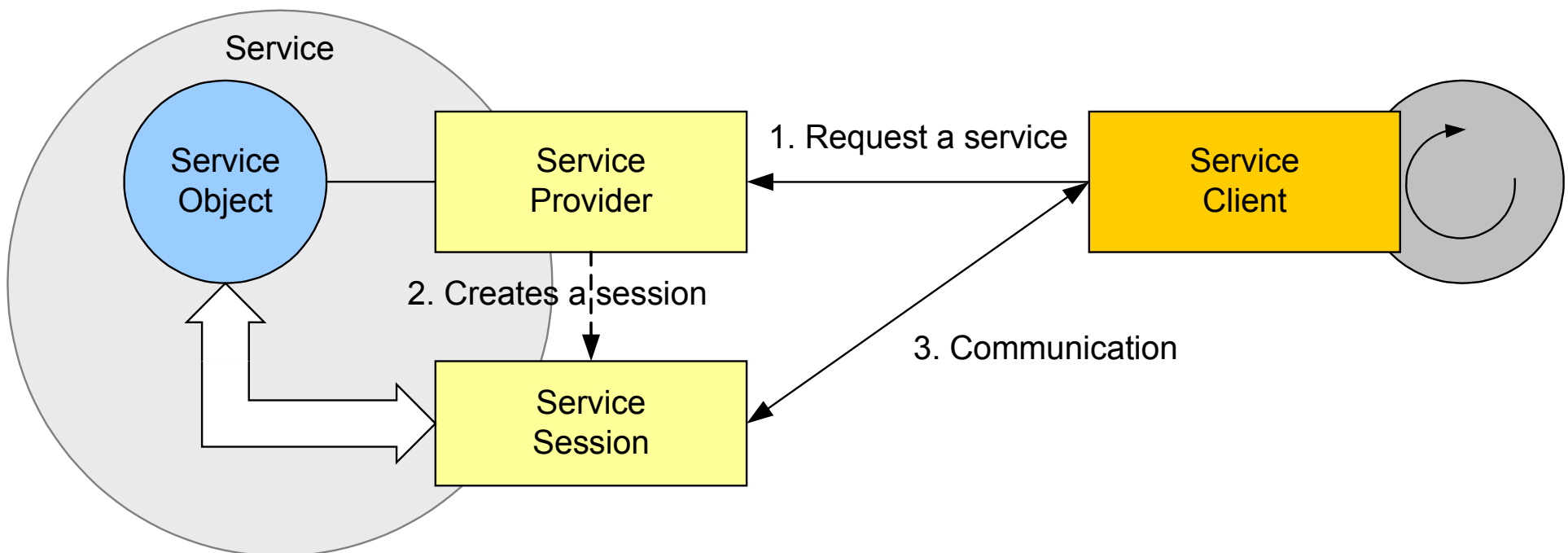
- Interface IPort
 - **open**: connects the port to a communication channel
 - **close**: disconnects the port from the channel
 - **peek**: gets a message from the channel without removing it
 - **read** and **readBlocked**: consumes a message from the channel
 - **write**: puts a message to the channel
- InputPort, OutputPort and InOutPort classes implement asynchronous port behavior

Services and Public Ports (I)

- Service: The service is a special entity providing a functionality or information to the world
 - one way communication: reading or writing only
 - two way communication: exchanging information, querying, interaction
- Public ports are accessible from outside (external systems)
 - Service Provider: The service provider port is used to provide services to entities
 - Service Session: Once connected to the service provider each entity is treated in a private session.
 - Service Client: Port on the client side

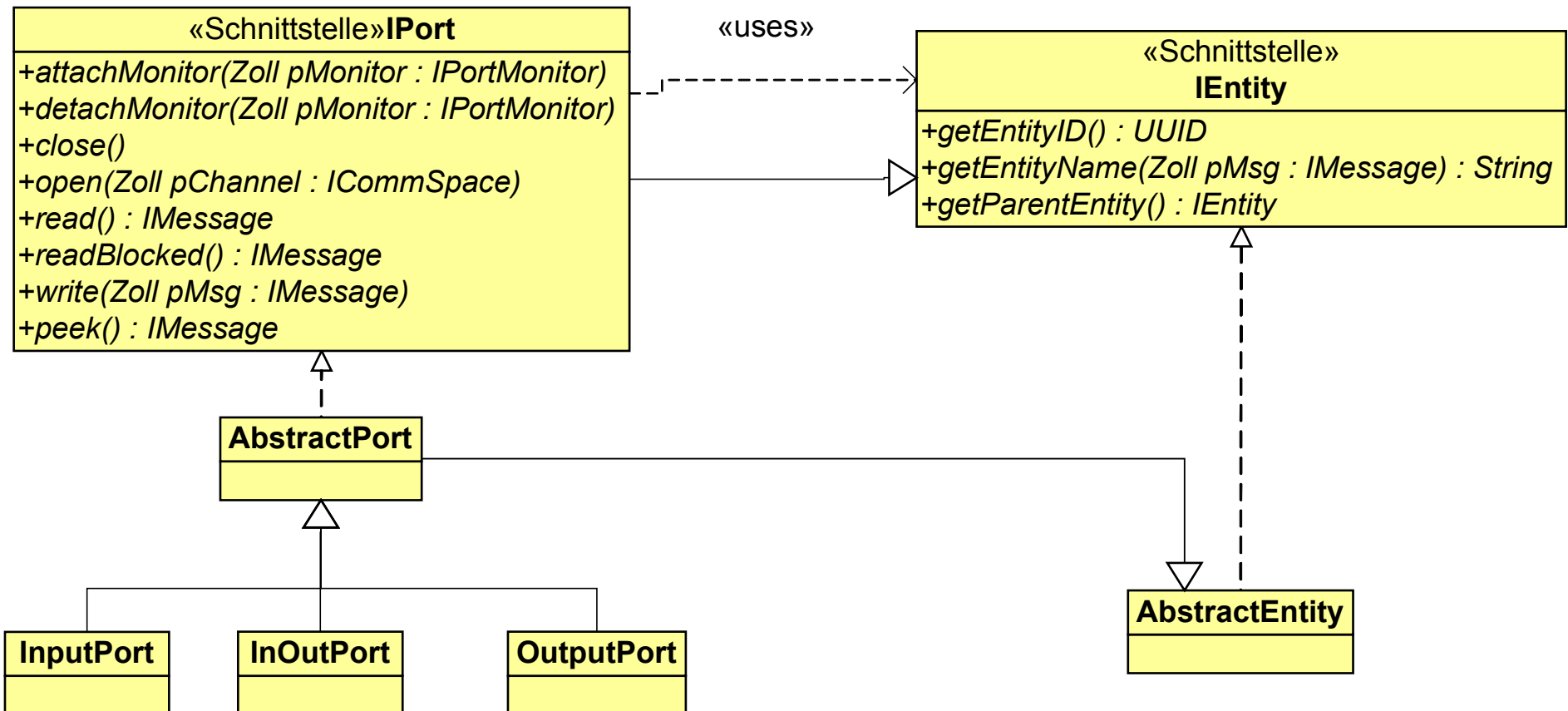
Services and Public Ports (II)

- 1.) requesting the access to service through public port
- 2.) if client is accepted a private session is created.
- 3.) the private session manages the access between client and service



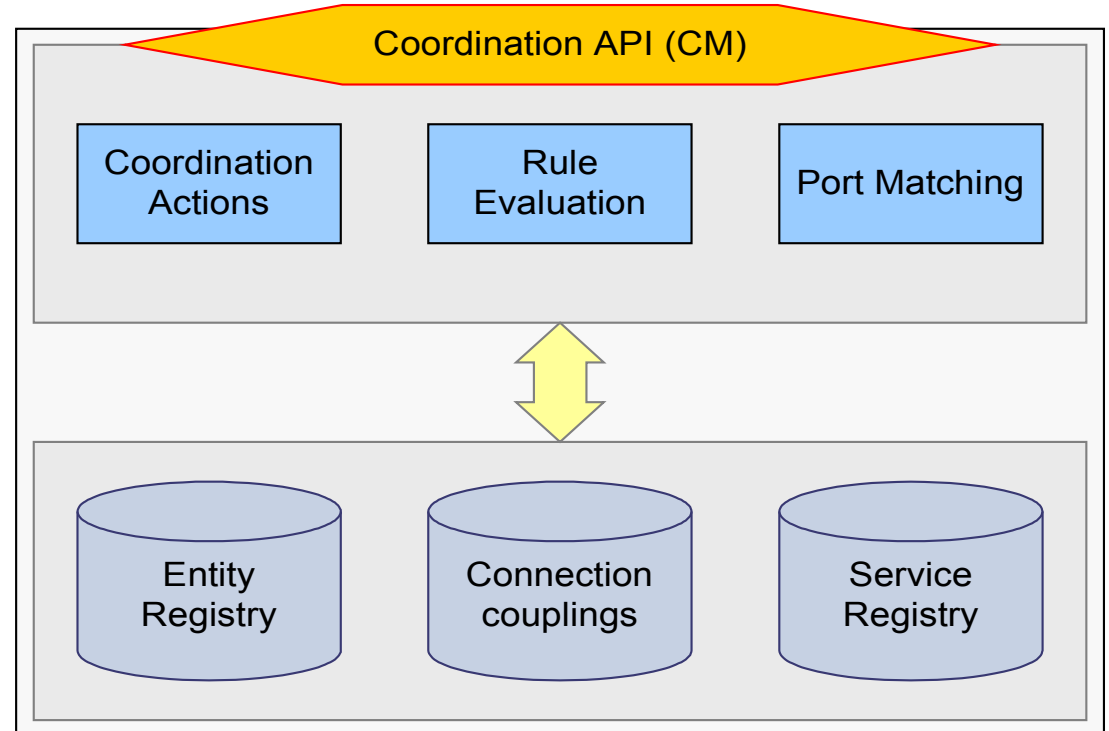
Ports: Classdiagram

- Implementation



Coordination Component (I)

- API for explicit coordination
 - port matching algorithm
 - rule evaluation
 - set of coordination actions



Coordination Component (II)

- The Coordination Manager is used to
 - register all available entities and ports
 - maintain the current connected port (port couplings)
 - clean and remove wasted channels
- The entity space observer is able to detect situation changes of entities within a system
 - evaluates the situation against rules
 - takes actions like
 - creating new channels between entities
 - disconnecting existing channels
 - providing or hiding services

Message Passing

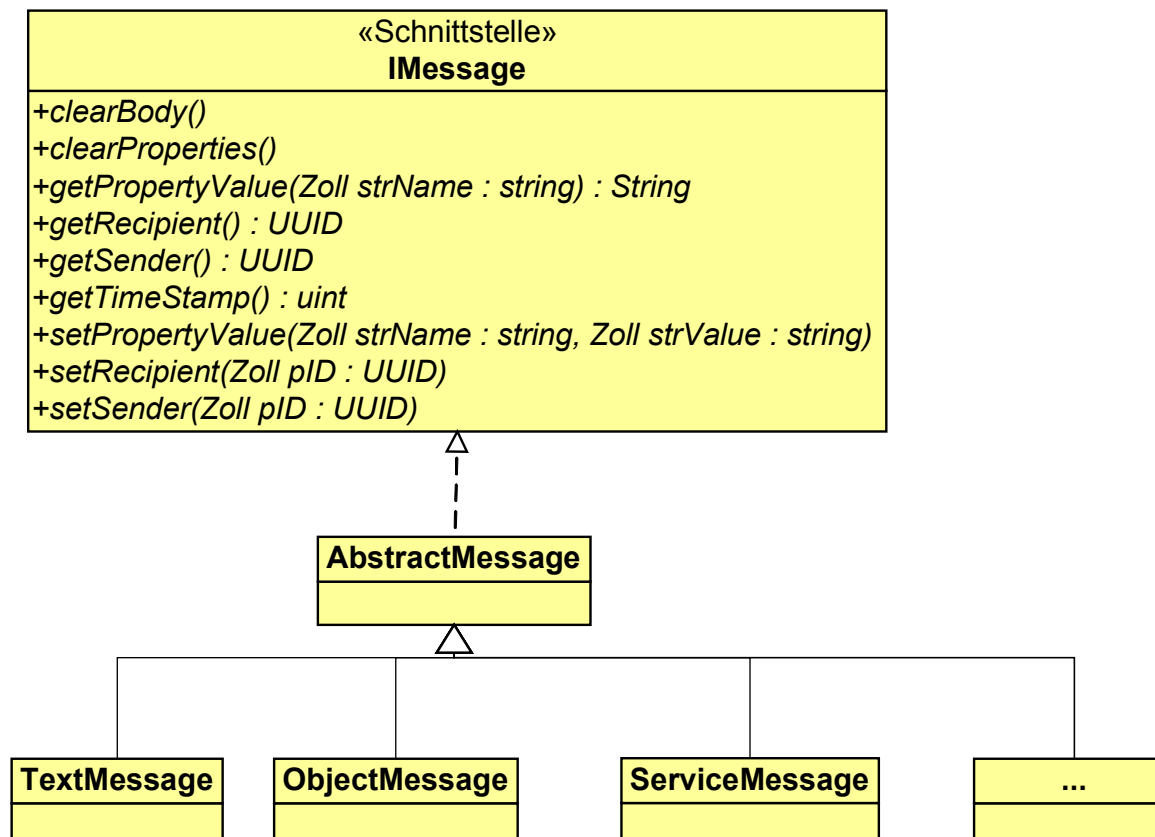
- The communication between entities is realized with message passing technique
 - Transparent for local and remote communication
 - sending through permanent channel
 - a channel is installed permanently between two entities
 - or sending through temporary channel (like e-mail)
 - a message is addresses for different recipients. To each recipient a temporary channel is opened. As soon as the recipient takes the message the channel is closed.

Message Passing Protocol (I)

- Basic construct is of transmission is a message object
 - Header: contains Message ID, Sender ID, Receiver IDs, logical timestamp, properties
 - Body: each concrete message implementation implements its own body
 - TextMessage: String message
 - ObjectMessage: Streamed object
 - ...
- A message is sent from one sender to one or many recipients

Message Passing Protocol (II)

- On top of the message protocol, application specific protocols can be defined

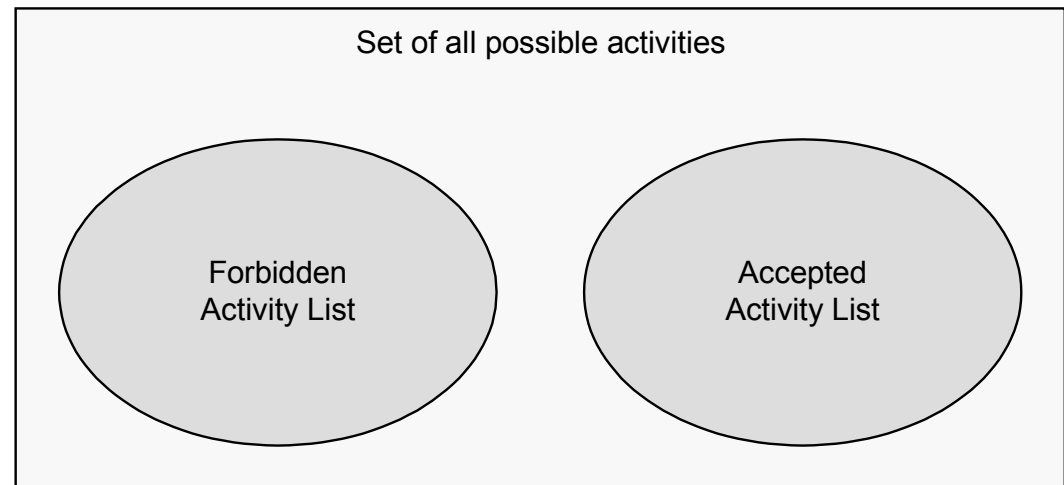


Rules

- Rules are used to
 - Check if an activity is allowed in a place
 - evaluate situations of entities
 - coordinate the interaction between two or more entities
- 3 different rules
 - activity rules
 - situation rules
 - social rules

Activity rules

- If an activity is recognized
 - the activity status is checked
 - the activity rules are attached to the place entity
- Activity lists
 - forbidden: not allowed activities.
 - accepted: appropriated activities.
 - everything else is negotiable.

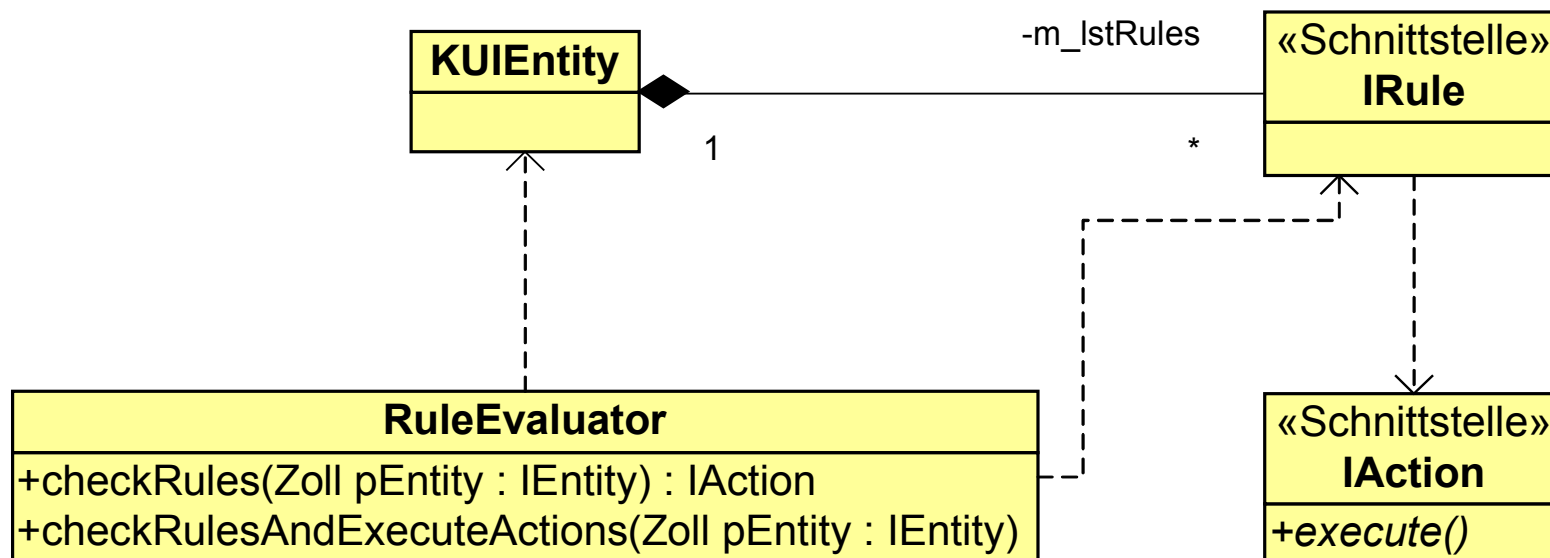


Situation rules

- Situation of each entity is treated in dedicated situation analysis
 - Situation Manager are attached to software observer
 - At each context, activity or structure change the situation is reevaluated
- The Situation Status is reported to the application
 - **Normal**: normal status
 - **Critical**: entity is in a critical situation
 - **Dangerous**: entity is in danger

Social Rules (I)

- Control the interaction between entities
 - sub entities inherit the rules from their parent
 - sub entities can overwrite and extend the rules given by the parent entities
 - after the rule evaluation a list of actions is returned



Social Rules (II)

- Rules inherit the interface IRule
 - **check()**: Checks if the rule is respected. If something has to be done the method returns an action object
 - **getName()**: returns the name of the rule. The name is needed to identify the rule
 - **isFinal()**: Final rules can not be overwritten by rules from child entities.

Merging Social Rules

- Merger rules using the rule-union operator \hat{U} :
 - R_1 and R_2 are set of rules from entity e_1 and e_2
(e_1 = parent of e_2)
 - if a rule exists in both sets the rule of the child entity e_2 is taken

$$R_1 \hat{U} R_2 = R_{xy}(R_1, R_2) \cup R_x(R_1, R_2) \cup R_y(R_1, R_2)$$

$$R_{xy}(R_1, R_2) = \{x \in R_1, y \in R_2 : id(x) \neq id(y)\}$$

$$R_x(R_1, R_2) = \{x \in R_1 : \exists y \in R_2 : id(x) = id(y) \wedge isFinal(x)\}$$

$$R_y(R_1, R_2) = \{y \in R_2 : \exists x \in R_1 : id(x) = id(y) \wedge \neg isFinal(x)\}$$